HIGH LEVEL COST EFFECTIVE MULTI-CLOUD DATA HOSTING SCHEME WITH HIGH AVAILABILITY

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ABSTRACT: Today most of the entrepreneur and companies are hosting their data on the cloud to reduce their cost as well as to increase their data availability. There are so many cloud vendors available in the market but there price and quality is major issue. Today we prefer the cloud which we can store most of the data and also looking for the price is cheaper compare than others. Most of the customers put all the data into a single cloud and then simply trust their luck this paper we can developed for a novel data hosting scheme (named CHARM) which combine two key functions desired. The first is selecting several suitable clouds and an appropriate redundancy strategy to store data with minimum monetary cost and 100 percent availability. The second is triggering a transition process to separate data according to the variations of data access pattern and pricing of clouds. We can calculate the performance of CHARM using both of the technique. The results show that compared with the major existing schemes; CHARM not only saves around 20 percent of monetary cost but also exhibits sound adaptability to data and price adjustments. Keyword: cloud, hosting, CHARAM

I. INTRODUCTION

There are so many cloud storage services are available in the market such as Drop box, CloudMe, and Seafile. These services provide users with a convenient and effective way to store and share data from anywhere on any device and at any time. User can store their information like documents, photos, videos and any other various file in cloud storagewhich are automatically synchronized across all the designated devices connected to the cloud in a timely manner. User can use their data from anywhere at any time with the help of mobile device availability. These features significantly simplify data management and consistency maintenance, and thus provide an ideal tool for data sharing and collaboration. Here we can have some comparison Cloud Me claims that over 200 million customers have stored more than 14 PB of data using their service, while Drop box has results more than 100 million users whocan store or update 1 billion files every day. Seafile enter into the market very late but still they are obtaining 10 million users just in its first two months. Their key features is of cloud storage services is data synchronization (sync) which automatically maps the changes in users' local file systems to the cloud via a series of network communications. There is one folder available in the storage where each and every file operation is noticed and synchronized by client software which is developed by service provider.synchronizing of a file means each and

every data to be sync like event, their index number, notification there status and also there acknowledgement. here in this paper we can mention traffic as synchronization of data traffic. day by day usage of internet is increasing and platform for online services also increasing like internet search, social networking and video streaming all of these data are to be distributed across multiple location for their better availability security purpose and reliability purpose. we can also improve smaller entrepreneur business with cloud computing platform like Amazon AWS. these service provide effective to access their data across different service location. most of the companies now a days offer distributed service to remove unnecessary burden. in this paper we can also design the structure implementation evaluation and how the system is deploy. we have to focus on two major criteria for better improvement. firstly we can focus that if a specified request is been generated then only specified respond will be given which complicates load balancing. second interactivity is also an important function in which application must respond within a specified time. according to these criteria hybrid architectures also implemented for entrepreneur application. majority of enterprise application are highly sensitive they require authentic key to access their data. some of the sensitive information like credit card number debit card number may be stored in the local to enterprise while other information can be stored in the cloud.

II. EXISTING SYSTEM

in the existing system the features

of data availability is measured by duplication of data or remove coding. in this paper we can design such an application which can meet different availability requirements. for replication replica can be stored in the several cloud so one data is removed so we can easily get another data from another server. cloud services are charges minimum

for their outgoing data bandwidth. for erasing coding data is encoded into n blocks and m coding blocks and all of these blocks are located into different cloud storage. in that case data availabiliy we can achieve with lower storage space.

DISADVANTAGES OF EXISTING SYSTEM

- Single Server availability
- Encryption process is not implemented
 - Does not provide any guarantee

III. PROPOSED SYSTEM

in this paper we can propose cost effective scheme name CHARM. it can be work like this data can be stored in

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multiple cloud with minimum cost and all time availabiliy. we can also combine two mechanism like duplication of data and reducing the code into a uniform manner to meet their special requirement.after the consideration of this scheme this structure will save around 20% in the overall cost.

ADVANTAGES

According to the duplicate mechanism the storage file size is small.

Their storage capacity depending upon their price. If the price high their storage automatically increase.

MODULE DESCRIPTION:

- Multi-cloud
- Data hosting
- Cloud Storage
- Server Module
- Consumer Module

MULTI-CLOUD

There are Lots of data centers are distributed around the world, and one region such as America, Asia, usually has several data centers belonging to the same or different cloud providers. in technically user can access data centers in a specified region. in some of the cases the access time of data centers is very low while some of the data centers is highly intolerable. CHARM chooses cloud for storing their data which meet their requirement. The capacity of storage mode does not impact on the performance of the service. This service is not a sensitive process we can also reduce the priority of operation and implemented them in batch wise when the proxy has major workload



DATA HOSTING

in this data hosting the whole model is located in the proxy there are four major component in CHARM.

Data hosting, storage mode switching, workload statistic and predictor. out of these four term workload statistic can collect and tackle all the access logs of data. Theterm also sends the statistic information to predictor which can give you the notification of SMS. Data hosting can store the data using duplication and erasing coding. Storage capacity can run in the background process.Predictor is used to predict the future access of different files. the time interval for prediction is one month after that we can access the file in the next month. There is lots of good algorithm for prediction. data hosing and SMS are two major functions in CHARM..



CLOUD STORAGE

Now a days cloud storage service is becoming the most popular. There are various reason behind that like importance of privacy, many cloud storage encryption schemes have been proposed to protect data from those who do not have access. All of these function can make cloud storage providers are safe and cannot be hacked in some of the cases they may reveal users data or have some access to their private data on the cloud, thus altogether circumventing storage encryption schemes. This paper can represent the design for new cloud storage encryption scheme that can prevent cloud storage providers to create convincing fake user secrets to protect user privacy. They cannot tell if obtained secrets are true or not, the cloud storage providers make sure thatuser privacy is firstly securely protected.Most of the proposed schemes assume cloud storage service providers or trusted third parties handling key management are trusted and cannot be hacked; however, in practice, some entities may intercept communications between users and cloud storage providers and then compel storage providers to release user secrets by using government power or other means. In this case, encrypted data are assumed to be known and storage providers are requested to release user secrets.We aimed to build an encryption scheme that could help cloud storage providers avoid this predicament. In our approach, we offer cloud storage providers means to create fake user secrets. Given such fake user secrets, outside coercers can only obtained forged data from a user's stored cipher text. Once coercers think the received secrets are real, they will be satisfied and more importantly cloud storage providers will not have revealed any real secrets. Therefore, user privacy is still protected. This concept comes from a special kind of encryption scheme called deniable encryption.

SERVER MODULE

This module is used to upload their files using some policy access. First of all they require public key for the file upload and then they require secret key for the file access. Using this secret key customer can upload their file.



CONSUMER MODULE

This function is used to help their client to search their file using their file id and name. if both of the data do not match then they do not get access to that file otherwise server can ask for their public key to access their file.



SCREENSHOTS



Owner Home:



View Services:

upload File to Cloud Server:





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IV. CONCLUSION

Cloud storage services are experimenting rapid development and this services are based on multi-cloud also become prevailing. In this paper we have designed a novel storage scheme CHARM which helps customer to distribute their data among the cloud very effectively and cheaper.CHARM makes fine-grained decisions about which storage mode to use and which clouds to place data in. The evaluation proves the efficiency of CHARM.

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